

In this lecture we look at the theory behind Simulated Annealing, and we will show that the method is guaranteed to converge - with probability 1 - for any starting value.

Simulated Annealing (§ 17.2)

- Start with an initial temperature T_0 .
- Start with an initial state of the system, $\underline{x}^{(0)}$.
- Generate a proposal to move the system into a new state $\underline{x}^{(1)}$.
- If the new state reduces the cost function, the proposal is accepted with probability 1.

$$E^{(1)} = E(\underline{x}^{(1)}), \quad E^{(0)} = E(\underline{x}^{(0)})$$

$$\Delta E = E^{(1)} - E^{(0)}$$

$$\mathbb{P}(\text{Accept } \underline{x}^{(0)} \rightarrow \underline{x}^{(1)}) = \begin{cases} 1 & \text{if } \Delta E < 0. \\ e^{-\Delta E / T_0} & \text{if } \Delta E > 0. \end{cases}$$

- After a certain number of steps like this, the temperature is reduced to T_1 .
- Continue thus ...
- Temperature is reduced in a systematic way according to an annealing schedule.

The algorithm satisfies detailed balance :

$$P_T(E^{(k)}) \mathbb{P}(\underline{x}^{(k)} \rightarrow \underline{x}^{(k+1)}) \\ = P_T(E^{(k+1)}) \mathbb{P}(\underline{x}^{(k+1)} \rightarrow \underline{x}^{(k)})$$

where $P_T(E) = \frac{e^{-\beta E}}{Z}$.

Detailed balance guarantees that all parts of the phase space are sampled.

Algorithm — § 17.3

Annealing Schedule — § 17.4

Aim: Justify the fact that

$$T_k \leq \frac{T_0}{\ln(k)}$$

is a good annealing schedule.

We start with an initial guess \underline{x}_0 . We generate a new guess (\underline{x}_1) — called a proposal — by drawing a vector from a normal distribution:

$$\underline{x}_1 \sim \mathcal{N}(\underline{x}_0, \sigma_k)$$

We take $\sigma_k = \sqrt{T_k}$ ($\sigma_k^2 = T_k$). Other references use $\sigma_k^2 = 2T_k$. Both choices

references use $\sigma_k^2 = 2T_k$. Both choices work. The choice $\sigma_k^2 = T_k$ is called Boltzmann annealing.

The probability that the new proposal is in a particular region R of phase space is:

$$P(\text{New proposal in } R) = \int_R p(\underline{x}) d^n \underline{x}$$

where $p(\underline{x})$ is the normal distribution

$$p(\underline{x}) = \frac{1}{(2\pi\sigma_k^2)^{n/2}} e^{-\|\underline{x} - \underline{x}_0\|^2 / 2\sigma_k^2}$$

As indicated previously, we propose the annealing schedule

$$T_k \leq \frac{T_0}{\ln k},$$

where T_0 is the initial temperature which is supposed to be "sufficiently large" (TBC).

We now look at the probability g_k that the new proposal is to be found in a small volume ΔV around the global min, and we denote this region as $R(\underline{x}^*, \Delta V)$. Hence:

$$g_k = P(\text{New proposal in } R(\underline{x}^*, \Delta V))$$

$$g_k = \mathbb{P}(\text{New proposal in } \mathcal{R}(\underline{x}_*, \Delta V))$$

$$= \int_{\mathcal{R}(\underline{x}_*, \Delta V)} p(\underline{x}) d^n \underline{x}$$

We have:

$$g_k = \int_{\mathcal{R}(\underline{x}_*, \Delta V)} \frac{1}{(2\pi\sigma_k^2)^{n/2}} e^{-\|\underline{x} - \underline{x}_*\|_2^2 / 2\sigma_k^2} d^n \underline{x}$$

$$= \left[\int_{x_{*1} - \Delta x/2}^{x_{*1} + \Delta x/2} \frac{1}{(2\pi\sigma_k^2)^{1/2}} e^{-(x_1 - x_{*1})^2 / 2\sigma_k^2} dx_1 \right] \times \dots$$

$$\dots \times \left[\int_{x_{*n} - \Delta x/2}^{x_{*n} + \Delta x/2} \frac{1}{(2\pi\sigma_k^2)^{1/2}} e^{-(x_n - x_{*n})^2 / 2\sigma_k^2} dx_n \right]$$

$$\approx \left[\frac{1}{(2\pi\sigma_k^2)^{1/2}} e^{-(x_{*1} - x_{*1})^2 / 2\sigma_k^2} \Delta x \right] \times \dots$$

$$\dots \times \left[\frac{1}{(2\pi\sigma_k^2)^{1/2}} e^{-(x_{*n} - x_{*n})^2 / 2\sigma_k^2} \Delta x \right]$$

$$\approx \frac{(\Delta x)^n}{(2\pi\sigma_k^2)^{n/2}} e^{-\|\underline{x}_* - \underline{x}_*\|_2^2 / 2\sigma_k^2}$$

$$= \Delta V e^{-\|\underline{x}_* - \underline{x}_*\|_2^2 / 2\sigma_k^2}$$

$$= \frac{\Delta V}{(2\pi\sigma_k^2)^{n/2}} e^{-\|\underline{x}_* - \underline{x}\|_2^2 / 2\sigma_k^2}.$$

Here, we have used the TRAPEZOIDAL RULE:

$$\int_{a-\Delta x/2}^{a+\Delta x/2} f(x) dx \approx f(a) \Delta x$$

Valid for smooth functions.

Summarizing where we are so far:

$$g_k = \frac{\Delta V}{(2\pi T_k)^{n/2}} e^{-\|\underline{x}_* - \underline{x}\|_2^2 / 2T_k}$$

We compute the probability that the system will NOT enter this region, and we will show that this probability goes to zero.

To show: $\prod_k (1 - g_k) \rightarrow 0$
 $\infty \quad k \rightarrow \infty.$

Take logs on both sides:

To show :

$$\log \prod (1 - g_n) \rightarrow -\infty \text{ as } k \rightarrow \infty.$$

Or :

$$\sum_k \log(1 - g_n) \rightarrow -\infty \text{ as } k \rightarrow \infty.$$

Assuming g_n is small, and Taylor-expand:

$$\log(1 - g_n) \approx -g_n.$$

Hence, it suffices to show:

$$\sum_k g_n \rightarrow +\infty \text{ as } k \rightarrow \infty.$$

Or

$$\sum_k \frac{1}{(2\pi T_k)^{1/2}} e^{-\|x_k - x_0\|_2^2 / 2T_k} \rightarrow \infty.$$

But : $T_k \leq \frac{T_0}{\ln(k)}$

Hence, it suffices to show:

$$\sum_k (\ln k)^{1/2} e^{-\ln k \|x_k - x_0\|_2^2 / 2T_0} \rightarrow \infty.$$

We choose T_0 large, such that:

$$\frac{\|x_n - x_0\|_2^2}{2T_0} \leq \Delta.$$

Then, it suffices to show:

$$\sum_k (\log k)^{n/2} e^{-\frac{(\log k) \|x_n - x_0\|_2^2}{2T_0}} \geq \sum_k (\log k)^{n/2} e^{-\log k}.$$

We look at the tail of the series, $k \geq k_0$:

$$\begin{aligned} & \sum_{k=k_0}^{\infty} (\log k)^{n/2} e^{-\frac{(\log k) \|x_n - x_0\|_2^2}{2T_0}} \\ & \geq \sum_{k=k_0}^{\infty} (\log k)^{n/2} e^{-\log k} \\ & = \sum_{k=k_0}^{\infty} (\log k)^{n/2} \cdot \frac{1}{k} \\ & \geq \sum_{k=k_0}^{\infty} \frac{1}{k} \end{aligned}$$

This is the divergent harmonic series.

..... $\geq \infty$.

Hence, $\sum_k (\log k)^{n/2} e^{-\frac{\|x_n - x_0\|_2^2}{2T_0} \log k} \rightarrow \infty$
 ∞ as $k \rightarrow \infty$.

Reasoning back up the chain,

$$\prod_{k=2}^n (1 - g_k) \rightarrow 0.$$

Probability of reaching the global min
(ΔV around the global min) goes to 1.

