

Week 4, Lectures 2-3

Motivation and Applications again (§ 1.2)

Example: A chemical company has the following set-up:

- 2 factories F_1 and F_2
- 12 retail outlets R_1, R_2, \dots, R_{12}

Furthermore,

- Each factory F_i can produce a_i tons of product per week (capacity)
- Each retail outlet has a demand b_j tons of product.

C_{ij} = cost of shipping 1 ton of product from factory i to outlet j .

X_{ij} = # tons shipped i to j .

Total cost:

$$f = \sum_{i=1}^2 \sum_{j=1}^{12} X_{ij} C_{ij}$$

Constraints

Capacity $\sum_{j=1}^{12} X_{ij} \leq a_i$

Demand $\sum_{i=1}^2 x_{ij} \begin{cases} = b_j \\ \geq b_j \end{cases}$

Positivity $x_{ij} \geq 0 \quad i=1,2; \quad j=1,\dots,12.$

We can reshape the unknowns x_{ij} into a ~~matrix~~ vector $\underline{x} \in \mathbb{R}^{24}$. The problem is now linear in \underline{x} . LINEAR PROGRAMMING.

Continuous versus discrete optimization

Often, instead of $\underline{x} \in \mathbb{R}^n$ we can have $\underline{x} \in \mathbb{N}^n$ or $\underline{x} \in \{0,1\}^n$.

Examples:

- Power grid n different types of power plant (wind, solar, nuclear, ...)

PP1, PP2, ..., PPn.

$$\underline{x} = (\dots, \underset{\substack{\uparrow \\ \in \mathbb{N}_0}}{x_j}, \dots, \dots)$$

$\underline{x} \in \mathbb{N}_0^n.$

- Putting a factory in city i :

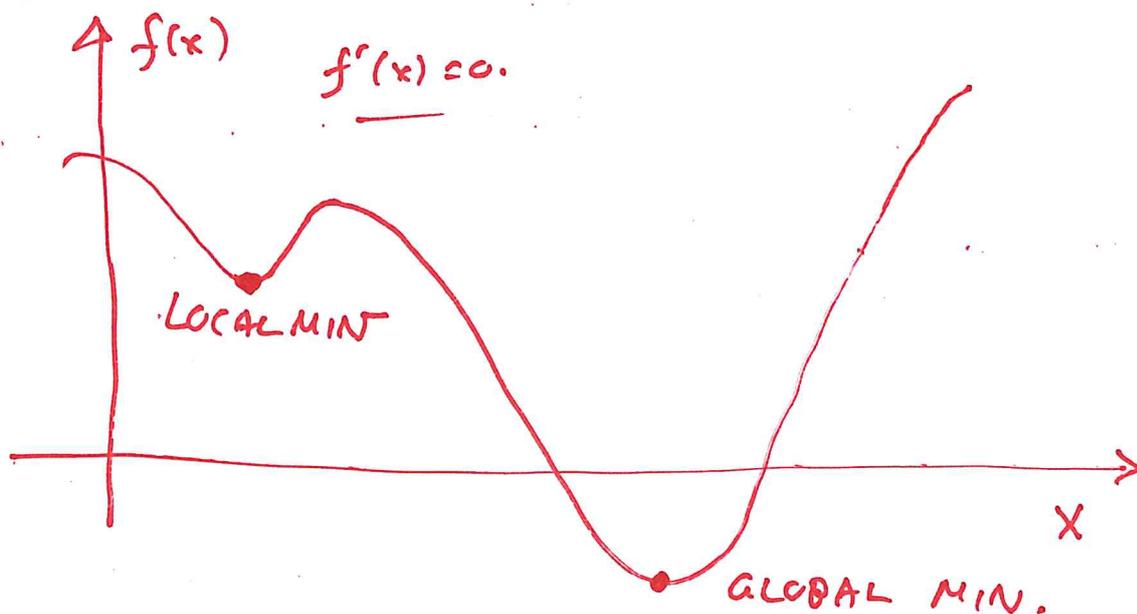
City 1, .., city n

$$Y(N) \begin{pmatrix} \phantom{0 \text{ or } 1} \\ \phantom{0 \text{ or } 1} \\ \phantom{0 \text{ or } 1} \\ \phantom{0 \text{ or } 1} \end{pmatrix}$$

$$\underline{x} \in \{0,1\}^n.$$

We focus in this module on continuous optimization.

Global versus local optimization.

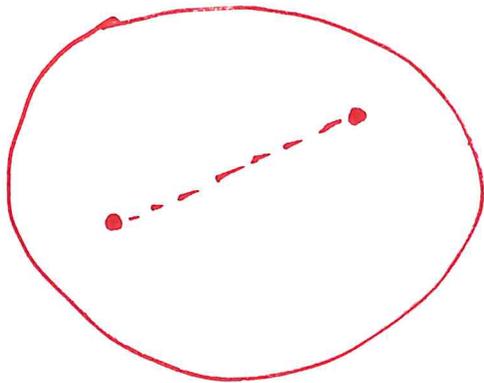


- Linesearch / Trust-Region / Newton will find a local min.
- Global optimizers (e.g. simulated annealing) will find the global min.

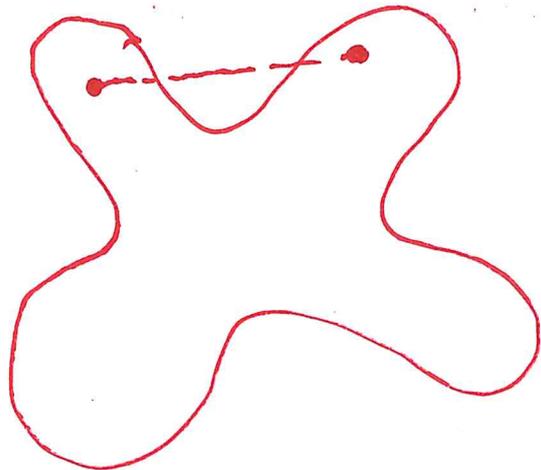
For convex functions there is a unique minimum.

§ 1.3 Convex sets

Defⁿ: Let $S \subset \mathbb{R}^n$. S is called convex if ~~a~~ straight-line segment joining any two points in S , is ~~entirely~~ entirely contained in S .



CONVEX .



NOT CONVEX .

S is convex, if, given any \underline{x} and \underline{y} in S , the line segment

$$\underline{x}(t) = \underline{x}t + (1-t)\underline{y} \quad t \in [0,1]$$

is contained entirely in S .

Examples :

Unit Ball

Exam
§1.3

Dot-product notation:

$$\langle \underline{x}, \underline{y} \rangle = \sum_{i=1}^n x_i y_i$$

$$\|\underline{x}\|_2^2 = \langle \underline{x}, \underline{x} \rangle = \sum_{i=1}^n x_i^2$$

Theorem 1.1 The unit ball

$$B = \{ \underline{x} \in \mathbb{R}^n \mid \|\underline{x}\|_2^2 \leq 1 \}$$

is a convex set.

Proof: ~~$\underline{x}(t)$~~ Let \underline{x} and $\underline{y} \in B$.

Take the line segment

$$\underline{x}(t) = \underline{x}t + (1-t)\underline{y}, \quad t \in [0,1]$$

Show: $\|\underline{x}(t)\|_2^2 \leq 1 \quad \forall t \in [0,1]$.

$$\begin{aligned} \|\underline{x}(t)\|_2^2 &= \langle \underline{x}(t), \underline{x}(t) \rangle \\ &= \langle \underline{x}t + (1-t)\underline{y}, \underline{x}t + (1-t)\underline{y} \rangle \\ &= t^2 \langle \underline{x}, \underline{x} \rangle + 2t(1-t) \langle \underline{x}, \underline{y} \rangle \\ &\quad + (1-t)^2 \langle \underline{y}, \underline{y} \rangle \\ &= t^2 \underbrace{\|\underline{x}\|_2^2}_{=1} + 2t(1-t) \langle \underline{x}, \underline{y} \rangle + (1-t)^2 \underbrace{\|\underline{y}\|_2^2}_{\leq 1} \\ &= t^2 + 2t(1-t) \langle \underline{x}, \underline{y} \rangle + (1-t)^2 \\ &\stackrel{\text{C.S.}}{\leq} t^2 + 2t(1-t) \underbrace{\|\underline{x}\|_2}_{\leq 1} \underbrace{\|\underline{y}\|_2}_{\leq 1} + (1-t)^2 \\ &= t^2 + 2t(1-t) + (1-t)^2 \\ &= \cancel{t^2} + \cancel{2t} - \cancel{2t^2} + 1 - \cancel{2t} + \cancel{t^2} \\ &= 1 \end{aligned}$$

$$\Rightarrow \|\underline{x}(t)\|_2^2 \leq 1 \quad \forall t \in [0,1].$$

Hence, $\underline{x}(t) \in B \quad \forall t \in [0,1]$.

So B is convex.

Polyhedra :

$$S = \{ \underline{x} \in \mathbb{R}^n \mid A\underline{x} = b, C\underline{x} \leq d \}$$

where A and C are matrices of appropriate dimension and b and \underline{d} are vectors.

$C\underline{x} \leq d$ means: $[C\underline{x}]_i \leq d_i$.

S is a convex set.

Convex Function :

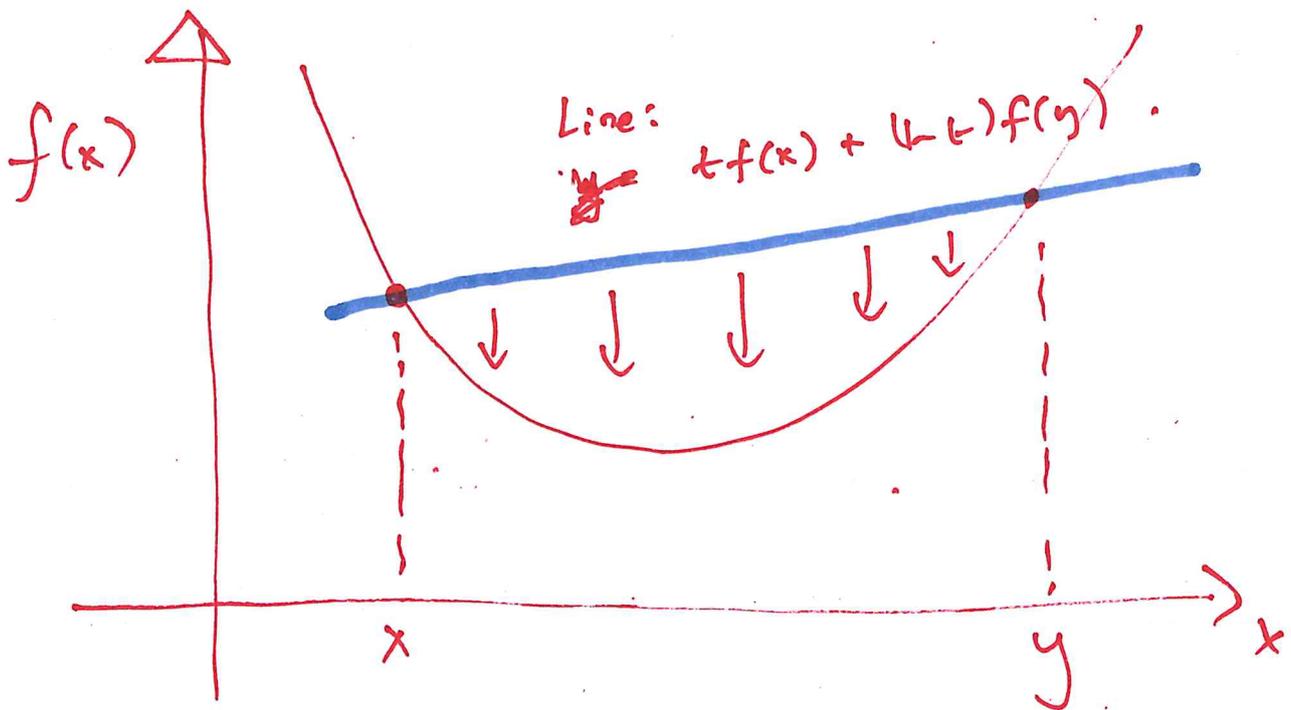
Let $f: S \rightarrow \mathbb{R}$, where $S \subset \mathbb{R}^n$.

The function f is convex if :

- S is a convex set
- The following relation holds :

$$f(t\underline{x} + (1-t)\underline{y}) \leq t f(\underline{x}) + (1-t)f(\underline{y})$$

for all $t \in [0,1]$ and all $\underline{x}, \underline{y} \in S$.



Examples:

- Let $S = \mathbb{R}^n$. Then the linear function

$$f(\underline{x}) = \langle \underline{b}, \underline{x} \rangle + \alpha$$

is convex. Here, \underline{b} is a const. vector and α is a constant. CHECK!

- Let $S = \mathbb{R}^n$. Then the following quadratic function is convex:

$$f(\underline{x}) = \langle \underline{x}, H\underline{x} \rangle \quad (1)$$

Here, H is symmetric semi-definite.

Theorem: The function in (1) is convex.

Proof: Let \underline{x} and \underline{y} be any two points in \mathbb{R}^n .

Look at:

$$\begin{aligned} & f(t\underline{x} + (1-t)y) \\ &= \langle t\underline{x} + (1-t)y, H[t\underline{x} + (1-t)y] \rangle \\ &= t^2 \langle \underline{x}, H\underline{x} \rangle + t(1-t) \langle \underline{x}, H\underline{y} \rangle \\ &\quad + (1-t)t \langle \underline{y}, H\underline{x} \rangle + \\ &\quad + (1-t)^2 \langle \underline{y}, H\underline{y} \rangle. \end{aligned}$$

Look at: $\langle \underline{y}, H\underline{x} \rangle = \langle H\underline{x}, \underline{y} \rangle$

$$\begin{aligned} &= \langle \underline{x}, H^T \underline{y} \rangle \\ &\stackrel{\text{Symmetric}}{=} \langle \underline{x}, H\underline{y} \rangle. \end{aligned}$$

Hence:

$$\begin{aligned} & f(t\underline{x} + (1-t)y) \\ &= t^2 \langle \underline{x}, H\underline{x} \rangle + 2t(1-t) \langle \underline{x}, H\underline{y} \rangle \\ &\quad + (1-t)^2 \langle \underline{y}, H\underline{y} \rangle \quad (*) \end{aligned}$$

Also, look at:

$$\begin{aligned} & t f(\underline{x}) + (1-t) f(\underline{y}) \\ &= t \langle \underline{x}, H\underline{x} \rangle + (1-t) \langle \underline{y}, H\underline{y} \rangle \quad (***) \end{aligned}$$

Take $(*) - (**)$:

$$\rightarrow t^2 - t = t(t-1)$$

$$t^2 \langle \underline{x}, H\underline{x} \rangle + 2t(1-t) \langle \underline{x}, H\underline{y} \rangle + (1-t)^2 \langle \underline{y}, H\underline{y} \rangle$$

$$- t \langle \underline{x}, H\underline{x} \rangle - (1-t) \langle \underline{y}, H\underline{y} \rangle$$

$$= (t^2 - t) \langle \underline{x}, H\underline{x} \rangle + 2t(1-t) \langle \underline{x}, H\underline{y} \rangle + [(1-t)^2 - (1-t)] \langle \underline{y}, H\underline{y} \rangle$$

$$= -t(1-t) \langle \underline{x}, H\underline{x} \rangle + 2t(1-t) \langle \underline{x}, H\underline{y} \rangle + (1-t) [1-t-1] \langle \underline{y}, H\underline{y} \rangle$$

$$= -t(1-t) \left[\langle \underline{x}, H\underline{x} \rangle - 2 \langle \underline{x}, H\underline{y} \rangle + \langle \underline{y}, H\underline{y} \rangle \right]$$

$$= \underbrace{-t(1-t)}_{\text{Pos. or zero}} \left[\underbrace{\langle (\underline{x}-\underline{y}), H(\underline{x}-\underline{y}) \rangle}_{\text{pos. or zero.}} \right]$$

$$\underbrace{\hspace{10em}}_{\text{neg. or zero}}$$

$$\leq 0.$$

Hence, $(*) - (**) \leq 0$, or
 $(*) \leq (**)$

Hence,

$$f(t\underline{x} + (1-t)\underline{y}) \leq t f(\underline{x}) + (1-t)f(\underline{y}) \quad \forall t \in [0,1].$$

Hence, f is convex. \square

Tuesday's lecture will be a recorded online lecture.

OP: $\min_{\underline{x} \in S} f(\underline{x}), \quad S \subset \mathbb{R}^n.$

Write the solution as \underline{x}_* :

$$\underline{x}_* = \arg \min_{\underline{x} \in S} f(\underline{x}). \quad (2)$$

Definition: \underline{x}_* is a global minimizer if

$$f(\underline{x}_*) \leq f(\underline{y}) \quad \forall \underline{y} \in S.$$

\underline{x}_* is a local minimizer if there exists a neighbourhood $\mathcal{N} \subset S$ such that $\underline{x}_* \in \mathcal{N}$, and such that

$$f(\underline{x}_*) \leq f(y) \quad \forall y \in \mathcal{N}.$$

Here, a neighbourhood means a non-empty open set.

\underline{x}_* is a strict local minimizer if there exists a neighbourhood $\mathcal{N} \subset S$ such that $\underline{x}_* \in \mathcal{N}$, and such that

$$f(\underline{x}_*) < f(y) \quad \text{for all } y \neq \underline{x}_* \in \mathcal{N}.$$

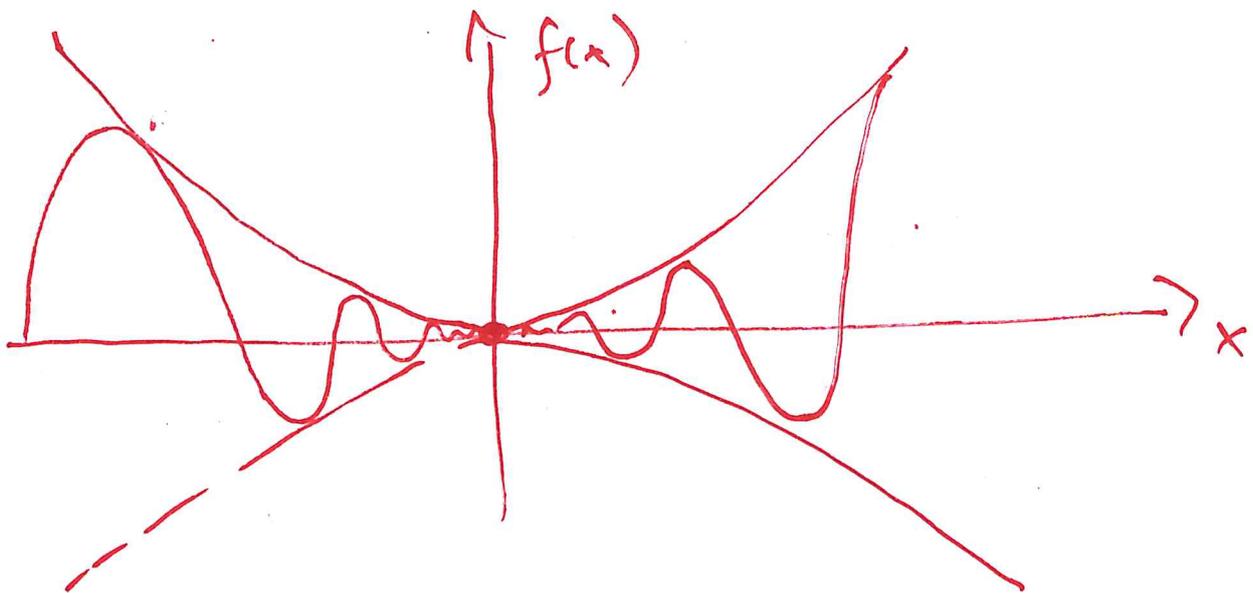
Examples:

$f(x) = 1$, $x \in \mathbb{R}^n$. Then every x is a local minimizer. But if I take $n = 1$, and $f(x) = (x-2)^4$, then $x_* = 2$ is a strict local minimizer.

\underline{x}_* is an isolated local minimizer if there exists a neighbourhood \mathcal{N} of \underline{x}_* such that \underline{x}_* is the only local minimizer in \mathcal{N} .

Pathological example:

$$f(x) = \begin{cases} x^4 \cos(1/x) + 2x^4, & x \neq 0. \\ 0 & x = 0. \end{cases}$$



There will always be a local min. ~~arbitrarily~~ arbitrarily close to the local min at $x=0$, so $x_0 = 0$ is not isolated.

§ 2.2. Necessary conditions for optimality when f is differentiable ($n=1$).

Taylor's Remainder Theorem: Let $f: \mathbb{R} \rightarrow \mathbb{R}$ be continuous with continuous first derivative. Assume $h > 0$. Then there exists $\eta \in (0, h)$ such that

$$f(x+h) = f(x) + f'(x+\eta)h.$$

Furthermore, if f is twice continuously differentiable, then there exists $\eta \in (0, h)$ such that:

$$f(x+h) = f(x) + f'(x)h + \frac{1}{2} f''(x+\eta)h^2.$$

$$f(x+h) \approx f(x) + f'(x)h + \frac{1}{2} f''(x)h^2.$$

Theorem 2.2 (First-order necessary condition). Let x_* solve the OP (2).

Then $f'(x_*) = 0$.

Proof by contradiction. Suppose that $f'(x_*) \neq 0$.

Look at $f'(x_*) < 0$. Then by continuity, there exists a $\delta > 0$ such that

$$f'(x) < 0 \text{ for all } x \in \underbrace{(x_* - \delta, x_* + \delta)}_I.$$

Look at $x = x_* + h$, where $|h| < \frac{1}{2}\delta$.

Apply Taylor's Remainder Theorem:

$$f(x) = f(x_*) + \overbrace{f'(x_* + \eta)h}^{\leftarrow x_* + h},$$

where $|\eta| < |h| < \delta$.

Since $x_* + \eta \in I$, $f'(x_* + \eta) < 0$.

Take $h = -k f'(x_* + \eta)$, $k > 0$.

Hence,

$$f(x) = f(x_*) - k |f'(x_*)|^2$$

Hence,

$$f(x) < f(x_*).$$

Contradiction, since x_* is the local minimizer.

The calculation for $f'(x_*) > 0$ is the same — a contradiction arises in both cases, hence:

$$f'(x_*) = 0.$$

Remark: The result for $n > 1$ is analogous:

$$\nabla f(x_*) = 0.$$

Theorem: (Second-order necessary condition)

Let x_* solve the OP (2), with $n = 1$.

Then $f''(x_*) \geq 0$.

Proof by contradiction. Assume that $f''(x_*) < 0$.

By continuity, there exists a $\delta > 0$ such that

$$f''(x) < 0 \text{ for all } x \in \underline{(x_* - \delta, x_* + \delta)}.$$

Consider $x = x_* + h$, with $|h| < \delta$. \bar{I}

$$f(x) = f(x_*) + \cancel{f'(x_*)h} + \frac{1}{2} f''(x_* + \eta) h^2,$$

where $|\eta| < |h| < \delta$, by Taylor's Remainder Theorem.

$$\Rightarrow f(x) = f(x_*) + \underbrace{\frac{1}{2} f''(x_* + \eta)}_{\text{neg.}} h^2.$$

$$\Rightarrow f(x) < f(x_*) \quad \text{Contradiction.}$$

$$\text{Hence, } \boxed{f''(x_*) \geq 0} \quad \square$$

Remark: In n -dimensions, the second-order necessary condition is that the Hessian matrix,

$$H_{ij} = \left(\frac{\partial^2 f}{\partial x_i \partial x_j} \right)_{x_*}$$

is positive-semi-definite.

$$\text{Reminder: } \langle \underline{\xi}, H \underline{\xi} \rangle \geq 0$$

for all $\underline{\xi} \in \mathbb{R}^n$.

Theorem: Second-order sufficient condition: Suppose that f is twice differentiable with continuous second derivative. If $f'(x_*) = 0$ and $f''(x_*) > 0$ (strict), then x_* is a strict local minimizer of f .

Proof: Idea the same as before (p.14).

Extension to n -dimensions:

Theorem 2.7 Suppose that f is twice differentiable with continuous second derivative. If $\nabla f(x_*) = 0$ and if the Hessian matrix H is positive-definite, then x_* is a strict local minimizer.

Reminder: H is positive-definite if $\langle \xi, H\xi \rangle > 0 \quad \forall \xi \neq 0 \in \mathbb{R}^n$.

Remark: Insisting on H being ~~sym~~ strictly positive definite rules out saddle points and degenerate critical points ($f(x) = x^4$).

For convex functions the situation is much nicer:

Theorem: When f is convex, any local minimizer ~~of~~ x_* is a global minimizer. If f is differentiable, then any stationary point ($\nabla f = 0$) is a global minimizer.

EXAM

Part 2. Show that a local minimizer really is a global minimizer. Proof by contradiction. Suppose that \underline{x}_* is a local minimizer, but that there exists $\underline{y} \in S$ such that

$$f(\underline{y}) < f(\underline{x}_*).$$

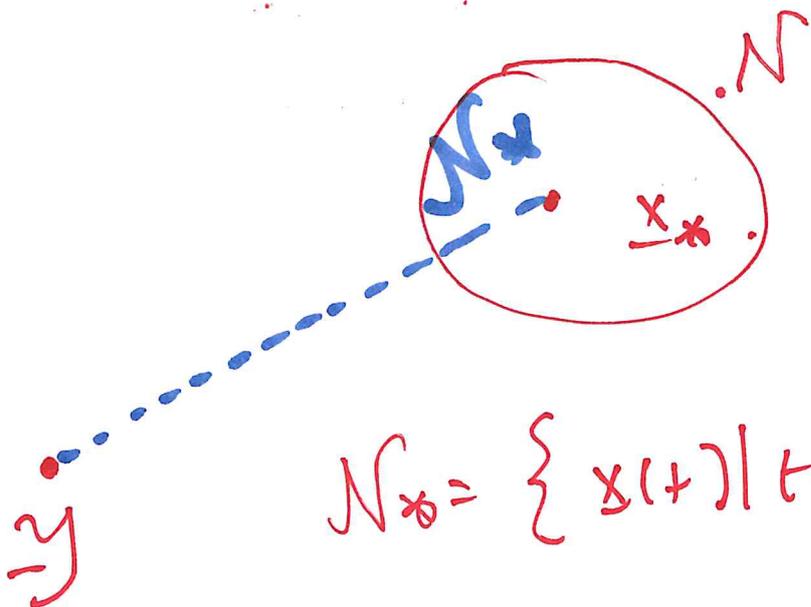
We construct the line segment

$$\underline{x}(t) = t\underline{y} + (1-t)\underline{x}_*, \quad t \in [0,1]$$

By convexity:

$$\begin{aligned} f(\underline{x}(t)) &\leq t f(\underline{y}) + (1-t) f(\underline{x}_*) \\ &\leq t f(\underline{x}_*) + (1-t) f(\underline{x}_*) \\ &= f(\underline{x}_*). \end{aligned}$$

$$\Rightarrow f(\underline{x}(t)) \leq f(\underline{x}_*).$$



$$N_{x_*} = \{ \underline{x}(t) \mid t \in [0,1] \} \cap N.$$

We have:

$$f(x(t)) \leq f(x_*) \quad \forall x(t) \in N_*$$

So no matter the size of N , there are points in N such that

$$f(\dots) \leq f(x_*).$$

This is a contradiction. Hence,

$$f(y) \geq f(x_*) \quad \forall y \in S.$$

Hence, x_* is a global min.

Part 2. Suppose that x_* is a critical point, $\nabla f(x_*) = 0$.

$$\lim_{t \rightarrow 0} \dots$$

$$\begin{aligned} 0 &= (y - x_*) \cdot \nabla f(x_*) \\ &= \frac{d}{dt} f(x_* + t(y - x_*)) \Big|_{t=0} \\ &= \lim_{t \downarrow 0} \frac{f(x_* + t(y - x_*)) - f(x_*)}{t} \\ &= \lim_{t \downarrow 0} \frac{f(ty + (1-t)x_*) - f(x_*)}{t} \\ &\leq \lim_{t \downarrow 0} \frac{tf(y) + \cancel{(1-t)f(x_*)} - \cancel{f(x_*)}}{t} \end{aligned}$$

$$\Rightarrow 0 \leq \lim_{t \downarrow 0} \frac{f(x_t) - f(x_*)}{t}$$

$$\Rightarrow 0 \leq f(y) - f(x_*)$$

$$\Rightarrow f(x_*) \leq f(y) \quad \forall y \in S.$$

Hence, x_* is a global min. \square