

On the Second Order Nonlinearity of a Cubic Maierana-McFarland Bent Function

Sumanta Sarkar¹ Sugata Gangopadhyay²

¹INRIA Paris-Rocquencourt, FRANCE

²Indian Institute of Technology, Roorkee, INDIA

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- ▶ Second order nonlinearity is an important cryptographic property.
- ▶ However, the best known algorithm to measure the second order nonlinearity of an n -variable Boolean function works for $n \leq 13$.
- ▶ Therefore, given a Boolean function, it is important to find a lower bound of its second order nonlinearity.
- ▶ As we know that bent functions have the maximum first order nonlinearity, it is interesting to check their second order nonlinearity.

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Our Contribution

- ▶ We study a new class of cubic Maiorana-McFarland bent functions which is based on a permutation constructed by Dobbertin (IEEE-IT 1999).
- ▶ First we show that this function can not have an affine derivative.
- ▶ Then we determine a lower bound of the second order nonlinearity of this function using Carlet's result (IEEE-IT 2008).

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Boolean Function

- ▶ Boolean function f is a mapping :

$$f : \{0, 1\}^n \rightarrow \{0, 1\}.$$

- ▶ This is an n -variable Boolean function.
- ▶ A Boolean function f can also be defined over the finite field \mathbb{F}_{2^n} :

$$f : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_2.$$

- ▶ Representing all $\alpha \in \mathbb{F}_{2^n}$ by the co-ordinates with respect to some basis of \mathbb{F}_{2^n} , the second representation gives the first one.

Reed-Muller Code

- ▶ The Reed-Muller code, $\mathcal{R}(r, n)$, of size 2^n and order r is the set of all n -variable Boolean functions of degree at most r .

r -th order nonlinearity

- ▶ Let f be an n -variable Boolean function.
- ▶ The r -th order nonlinearity ($nl_r(f)$) of f is the distance from f to the Reed-Muller code $\mathcal{R}(r, n)$.
- ▶ For $r = 1$, we simply denote it as nonlinearity.
- ▶ In this work, we are interested in $r = 2$, that is the second order nonlinearity.
- ▶ The r -th order nonlinearity is an important cryptographic property for block and stream ciphers.
- ▶ For example, there have been some notion of attack by using nonlinear approximations (r -degree Boolean function, $r > 1$) to f . To resist this attack the function needs to have high r -th order nonlinearity.

Maiorana McFarland Bent functions

- ▶ For even n , the maximum nonlinearity of an n -variable Boolean function is $2^{n-1} - 2^{\frac{n}{2}-1}$.
- ▶ For even n , the Boolean functions which possess this maximum nonlinearity are called bent functions.
- ▶ Maiorana McFarland is an important class of bent functions.
- ▶ Let $n = 2t$.
- ▶ The function $f : \mathbb{F}_{2^t} \times \mathbb{F}_{2^t} \rightarrow \mathbb{F}_2$ given by

$$f(x, y) = \text{Tr}_1^t(x\pi(y))$$

is a Maiorana-McFarland bent function where $\pi : \mathbb{F}_{2^t} \rightarrow \mathbb{F}_{2^t}$ is a permutation and

$$\text{Tr}_1^t(x) = x + x^2 + x^{2^2} + \dots + x^{2^{n-1}}.$$

Derivatives

- ▶ Let $f : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_2$.
- ▶ The derivative of f with respect to $a \in \mathbb{F}_{2^n}$, is denoted by $D_a f$ and is the Boolean function defined by

$$D_a f(x) = f(x) + f(x + a)$$

for all $x \in \mathbb{F}_{2^n}$.

Walsh Spectrum

- ▶ The Walsh transform of a Boolean function $f : \mathbb{F}_{2^n} \mapsto \mathbb{F}_2$ at $\lambda \in \mathbb{F}_{2^n}$ is defined as follows:

$$W_f(\lambda) = \sum_{x \in \mathbb{F}_{2^n}} (-1)^{f(x) + \text{Tr}(\lambda x)}.$$

- ▶ Walsh spectrum is the set $\{W_f(\lambda) : \lambda \in \mathbb{F}_{2^n}\}$.

Walsh Spectrum of quadratic Boolean functions

- ▶ The bilinear form associated to f is defined by

$$B(x, y) = f(0) + f(x) + f(y) + f(x + y).$$

- ▶ The kernel of $B(x, y)$ is the subspace defined by

$$\mathcal{E}_f = \{x \in \mathbb{F}_{2^n} : B(x, y) = 0 \text{ for all } y \in \mathbb{F}_{2^n}\}.$$

- ▶ For a quadratic Boolean function f , the kernel \mathcal{E}_f , is given by

$$\mathcal{E}_f = \{a \in \mathbb{F}_{2^n} \mid D_a f = \text{constant}\}.$$

Walsh Spectrum of quadratic Boolean functions (Continued)

- ▶ **Lemma 1:**(Macwilliams and Sloane; Canteaut, Charpin, Kyureghyan FFA 2008)

If $f : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_2$ is a quadratic Boolean function and $B(x, y)$ is the quadratic form associated to it, then the Walsh Spectrum of f depends only on the dimension, k , of the kernel, \mathcal{E}_f , of $B(x, y)$. The weight distribution of the Walsh spectrum of f is:

$W_f(\lambda)$	number of λ
0	$2^n - 2^{n-k}$
$2^{(n+k)/2}$	$2^{n-k-1} + (-1)^{f(0)}2^{(n-k-2)/2}$
$-2^{(n+k)/2}$	$2^{n-k-1} - (-1)^{f(0)}2^{(n-k-2)/2}$

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► **Proposition 1:** (Carlet IEEE-IT 2008)

Let f be any n -variable Boolean function and r be a positive integer smaller than n , then

$$nl_r(f) \geq 2^{n-1} - \frac{1}{2} \sqrt{2^{2n} - 2 \sum_{a \in \mathbb{F}_{2^n}} nl_{r-1}(D_a f)}.$$

► If exact values of $nl_{r-1}(D_a f)$ for all a are not known, but some lower bound is known, then we have the following corollary.

► **Corollary 1:** (Carlet IEEE-IT 2008)

Let f be any n -variable function and r be a positive integer smaller than n . Assume that for some nonnegative integers M and m , we have $nl_{r-1}(D_a f) \geq 2^{n-1} - M2^m$ for every nonzero $a \in \mathbb{F}_{2^n}$. Then

$$\begin{aligned} nl_r(f) &\geq 2^{n-1} - \frac{1}{2} \sqrt{(2^n - 1)M2^{m+1} + 2^n} \\ &\approx 2^{n-1} - \sqrt{M2^{\frac{n+m-1}{2}}}. \end{aligned}$$

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- ▶ Let $n = 2t$, where $t = 2m + 1$ and $m \geq 2$.
- ▶ We define the cubic Maiorana-McFarland function $\phi_n : \mathbb{F}_{2^t} \times \mathbb{F}_{2^t} \rightarrow \mathbb{F}_2$ as $\phi_n(x, y) = \text{Tr}_1^t(x(y^{2^{m+1}+1} + y^3 + y))$, where $y \mapsto y^{2^{m+1}+1} + y^3 + y$ is a permutation over \mathbb{F}_{2^t} (Dobbertin IEEE-IT 1999).

Derivatives of ϕ_n

Theorem 1: The function ϕ_n does not possess any derivative in $\mathcal{R}(1, n)$.

Lower bound of second order nonlinearity of ϕ_n

Theorem 2: The lower bound of the second order nonlinearity of ϕ_n is given as

$$\begin{aligned}nl_2(\phi_n) &\geq 2^{n-1} - \frac{1}{2}\sqrt{(2^n - 1)2^{\frac{n+t}{2}+3} + 2^n} \\ &\approx 2^{n-1} - 2^{\frac{7n+4}{8}}.\end{aligned}$$

Outline of the proof of Theorem 2

- ▶ Let $a, b \in \mathbb{F}_{2^t}$.
- ▶ Let $k(a, b)$ denote the dimension of the subspace

$$\mathcal{E}_{\phi_n} = \{(c, d) \in \mathbb{F}_{2^t} \times \mathbb{F}_{2^t} \mid D_{(c,d)} D_{(a,b)}(\phi_n) = \text{constant}\}.$$

- ▶ $k(a, b) = \begin{cases} t + i, & 0 \leq i \leq 4 \text{ when } b = 0 \\ r + j, & 0 \leq j \leq 2, 0 \leq r \leq 2 \text{ when } b \neq 0 \end{cases}$

Outline of the proof of Theorem 2 (Continued)

- ▶ $D_{(a,b)}(\phi_n)$ is always quadratic (by Theorem 1) for $(a, b) \neq (0, 0)$.
- ▶ By Lemma 1,

$$nl(D_{a,b}(\phi_n)) = 2^{n-1} - 2^{\frac{n+k(a,b)}{2}}.$$

- ▶ since $i \leq 4$,

$$nl(D_{a,b}(\phi_n)) \geq 2^{n-1} - 2^{\frac{n+t+4}{2}}.$$

Outline of the proof of Theorem 2 (Continued)

- ▶ Comparing with Corollary 1, we get $M = 1$ and $m = \frac{n+t}{2} + 2$.
- ▶ This gives

$$\begin{aligned}nl_2(\phi_n) &\geq 2^{n-1} - \frac{1}{2}\sqrt{(2^n - 1)2^{\frac{n+t}{2}+3} + 2^n} \\ &\approx 2^{n-1} - 2^{\frac{7n+4}{8}}.\end{aligned}$$

Better lower bound than the general bound

- ▶ The general lower bound of the second order nonlinearity (Carlet IEEE-IT 2008) of an n -variable cubic Boolean function which does not have any derivative in $\mathcal{R}(1, n)$ is

$$2^{n-1} - 2^{n-\frac{3}{2}}.$$

- ▶ $2^{n-1} - 2^{\frac{7n+4}{8}} > 2^{n-1} - 2^{n-\frac{3}{2}}$, for all $n > 16$.

Conclusions and further research

- ▶ We have identified a class of Majorana McFarland bent functions which do not have any affine derivative.
- ▶ We have studied the second order nonlinearity of these functions.
- ▶ Next step is to find a better lower bound of second order nonlinearity of this class of functions for which we need a new strategy.

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THANK YOU