

# Multiplicative Order of Gauss Periods

Omran Ahmadi<sup>1</sup> Igor Shparlinski<sup>2</sup> Jose Felipe Voloch<sup>3</sup>

<sup>1</sup>CSI, UCD

<sup>2</sup>Macquarie University

<sup>3</sup>University of Texas at Austin

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# 1, Notations

## Some Notations

- $\mathbb{F}_q$  is the finite field with  $q$  elements for a prime power  $q$ .
- $\mathbb{F}_{q^n}$  is the degree  $n$  extension of  $\mathbb{F}_q$ .
- Generators of  $\mathbb{F}_q^*$  are called **primitive** elements.

## Open Question

*Find an efficient algorithm for constructing primitive elements in finite fields.*

- An algorithm is efficient if its running time is  $(\log q^n)^{O(1)}$  arithmetic operations in  $\mathbb{F}_{q^n}$ .
- In many applications (Diffie-Hellman key establishment, pseudorandom bit generations, . . . ) a primitive element is needed.

## 3, Main Strategies

- 1 Find a small subset  $S \subset \mathbb{F}_q$  containing a primitive element.  
(Distribution result) (quite efficient assuming GRH)
- 2 Test the elements of  $S$  to find a primitive element.

## 4, Testing primitiveness

- (Naive) Compute all the powers of  $\alpha \in \mathbb{F}_q$ .
- $\alpha$  is primitive iff  $\alpha^{(q-1)/d} \neq 1$  for every prime  $d|q-1$ .
- (Bottleneck) Factorization (Subexponential time algorithm).
- Running time of number field sieve  
 $O(\exp((c + o(1))(\ln q)^{1/3}(\ln \ln q)^{2/3}))$  bit operations. (Best asymptotic running time)

## 5, Relaxation of Primitive Element Problem

- In many application we need an element of large order.
- Given  $\mathbb{F}_q$ , it suffices to construct primitive element for  $\mathbb{F}_{q^n}$  where  $n$  is in the some large subset of the positive integers.

## 6, Main Idea

- An element of  $\mathbb{F}_{q^n}$  which is not in any subfield has minimal polynomial of degree  $n$  over  $\mathbb{F}_q$ .
- If  $f(\alpha) = 0$ , then  $f(\alpha^q) = 0$ .
- Powers of an element and its degree are related through its minimal polynomial.
- AKS Deterministic Primality Testing Algorithm

## 7, Qi Cheng's Result

### Theorem (Q. Cheng, 2004)

*Let  $q$  be a fixed prime power and let  $N$  be a positive integer. Then in time polynomial in  $N$  an integer  $n \in [N, 2qN]$  and  $\alpha \in \mathbb{F}_{q^n}$  being of order at least  $5.8^{n/\log_q n}$  can be found.*

### Theorem (Q. Cheng, 2004)

*Let  $q$  be a fixed prime power and let  $N$  be a positive integer. Then in time polynomial in  $N$  an integer  $n \in [N, N + O(N^{0.77})]$  and  $\alpha \in \mathbb{F}_{q^n}$  being of order at least  $5.8^{\sqrt{n}}$  can be found.*



## 8, Main Idea and Proof

- Lemma: Let  $q$  be a prime power and let  $n|q - 1$ . If  $x^n - u \in \mathbb{F}_q[x]$  is an irreducible polynomial over  $\mathbb{F}_q$  and  $\alpha \in \mathbb{F}_{q^n}$  is one of its roots, Then for any  $a \in \mathbb{F}_q^*$ ,  $\alpha + a$  has order greater than  $5.8^n$ .
- Conjugates of  $\alpha$  are  $c_1\alpha, c_2\alpha, \dots, c_n\alpha$  where  $c_i$ 's are  $n$ -th roots of unity in  $\mathbb{F}_q$ .
- Conjugates of  $\alpha + a$  are  $c_i\alpha + a$ 's.

# 9, Proof

- Let  $e_i$ 's be positive integers such that  $\sum e_i \leq n - 1$ .
- Take the elements  $(\alpha + a)^{\sum e_i q^i} \in \mathbb{F}_{q^n}$ .
- $(\alpha + a)^{\sum e_i q^i} = \prod (\alpha + a)^{e_i q^i} = \prod (c_i \alpha + a)^{e_i}$ .
- If  $\prod (c_i \alpha + a)^{e_i} = \prod (c_i \alpha + a)^{f_i}$ , then the degree of  $\alpha \leq n$ . (a contradiction)
- $\alpha$  becomes a root of  $\prod (c_i x + a)^{e_i} - \prod (c_i x + a)^{f_i}$ .

# 10, Gauss Periods

## Definition

Let  $r = 2n + 1$  be a prime number coprime with  $q$  and  $\beta \in \mathbb{F}_{q^{2n}}$  be a primitive  $r$ -th root of unity. Then the element

$$\alpha = \beta + \beta^{-1} \in \mathbb{F}_{q^n} \quad (1)$$

is called a Gauss period of type  $(n, 2)$ .

## Theorem

If  $q$  is a primitive root modulo  $r$ , and  $\alpha$  is a Gauss period, then  $NB = \{\alpha, \alpha^q, \dots, \alpha^{q^{n-1}}\}$  is a normal basis for  $\mathbb{F}_{q^n}$  over  $\mathbb{F}_q$ . In this case, the minimal polynomial of  $\beta$  is of degree  $2n$ .

# 11, Orders of Gauss periods

Theorem (I. Shparlinski and J. von zur Gathen, 1998)

*Let  $p$  be the characteristic of  $\mathbb{F}_q$  and let  $q$  be a primitive root modulo a prime  $r = 2n + 1$  and  $\beta$  be a primitive  $r$ -th root of unity in  $\mathbb{F}_{q^{2n}}$ . Then the multiplicative order of  $\alpha = \beta + \beta^{-1}$  is at least*

$$2^{\sqrt{2n}}.$$

Theorem (ASV)

*With the assumptions as above the multiplicative order of  $\alpha$  is at least*

$$P(n - 1, p - 1).$$

$P(m, k)$  is the number of integer partitions of  $m$  where no part appears more than  $k$  times.

# 12, Orders of Gauss Periods

## Corollary (ASV)

*Assuming everything as the last theorem the multiplicative order of  $\alpha$  is at least*

$$\exp \left( \left( \pi \sqrt{\frac{2(p-1)}{3p}} + o(1) \right) \sqrt{n} \right),$$

as  $n \rightarrow \infty$ .

- When  $p > n$ , the order of  $\alpha$  is at least  $13\sqrt{n}$ .
- When  $p = 2$ , the order of  $\alpha$  is at least  $6.13\sqrt{n}$ .

# 13, Proof

- $\mathfrak{P} = \left\{ (u_1, \dots, u_{n-1}) \mid \sum_{j=1}^{n-1} u_j j = n-1, 0 \leq u_j \leq p-1 \right\}$ .
- Let  $q^{z_j} \equiv j \pmod{r}$ ,  $0 \leq z_j < r$ . ( $q$  is a primitive root modulo  $r$ )
- If  $(u_1, \dots, u_{n-1}), (v_1, \dots, v_{n-1}) \in \mathfrak{P}$  and  $(u_1, \dots, u_{n-1}) \neq (v_1, \dots, v_{n-1})$ , then  $\alpha^{\sum u_i q^{z_i}} \neq \alpha^{\sum v_i q^{z_i}}$ .
- $\alpha^{\sum u_i q^{z_i}} = \prod (\beta^{q^{z_i}} + \beta^{-q^{z_i}})^{u_i} = \prod (\beta^i + \beta^{-i})^{u_i} = \beta^{n-1} \prod (\beta^{2^i} + 1)^{u_i}$ .

# 14, Applying Polynomial ABC Theorem

- $\alpha \sum v_i q^{2i} = \beta^{n-1} \prod (\beta^{2i} + 1)^{v_i}$ .
- If  $\alpha \sum u_i q^{2i} = \alpha \sum v_i q^{2i}$ , then  $\prod (\beta^{2i} + 1)^{u_i} = \prod (\beta^{2i} + 1)^{v_i}$ .
- $\beta$  is of degree  $2n$ . (A contradiction)

## Theorem (Polynomial ABC theorem)

*Let  $A, B, C$  be nonzero polynomials over  $\mathbb{F}_q$  with  $A + B + C = 0$  and  $\gcd(A, B, C) = 1$ . If  $\deg A \geq \deg \text{rad } ABC$ , then  $A' = 0$ .*

- Using ABC we can beat GS bound but still unable to beat the partition bound.

# Outlook

- Outlook
  - Get similar bound for other types of Gauss periods.
  - Apply polynomial ABC theorem.